

イデアル格子で作る
効率の良い公開鍵暗号方式、他

草川 恵太 (東京工業大学)

結果は？

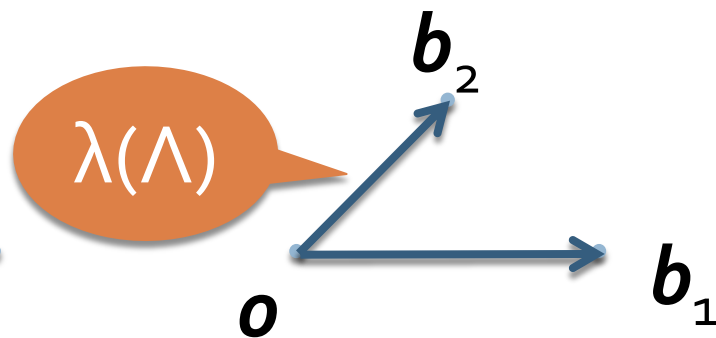
	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MR07, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

説明は？

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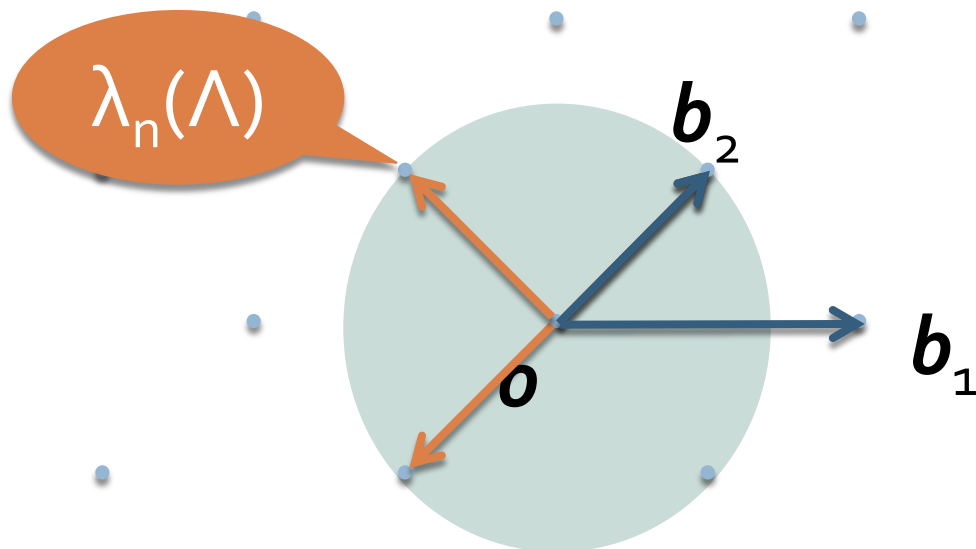
格子

- $\mathbf{B}=[\mathbf{b}_1, \dots, \mathbf{b}_n] \in \mathbb{Q}^{n \times n}$
- $L(\mathbf{B}) := \{\sum_i \alpha_i \mathbf{b}_i \mid \alpha_i \in \mathbb{Z} \text{ for all } i\} \subseteq \mathbb{Q}^n$
- $\lambda(\Lambda) : \Lambda$ 中の最短ベクトルの長さ



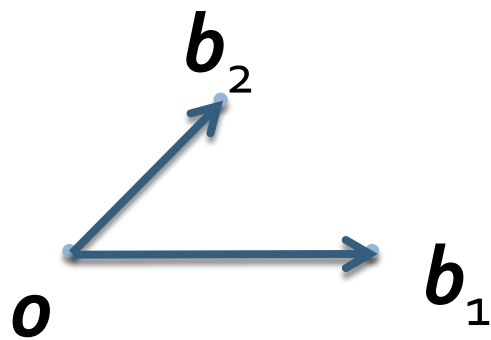
格子定数₂

- $\lambda_n(\Lambda)$: Λ 中の線形独立なベクトルの組の長さの最大値の最小値
- $\lambda_n(\Lambda) = \min_{S: \text{lin. ind. set in } \Lambda} \max_i \|s_i\|$
- $\lambda_n(\Lambda) = \min\{r: \dim(\text{span}(B_n(r) \cap \Lambda))=n\}$



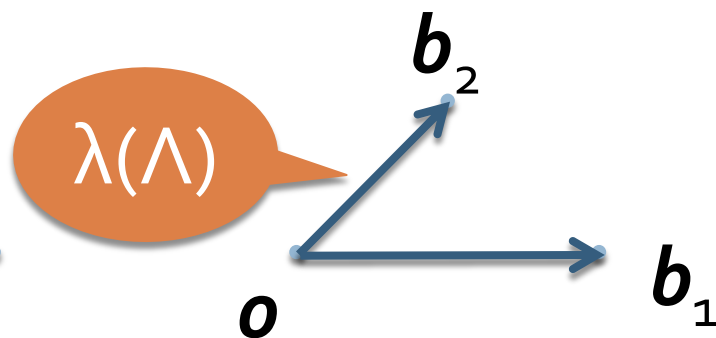
格子: 例

$$\mathbf{B} = \begin{pmatrix} 0 & 1 \\ 2 & 1 \end{pmatrix}$$



双对格子

- $B = [b_1, \dots, b_n] \in \mathbb{Q}^{n \times n}$
- $L(B) := \{\sum_i \alpha_i b_i \mid \alpha_i \in \mathbb{Z} \text{ for all } i\} \subseteq \mathbb{Q}^n$
- $\Lambda^* = \{x \mid \langle x, v \rangle \in \mathbb{Z} \text{ for all } v \text{ in } \Lambda\}$
- $\Lambda^* = L(B^{-T})$

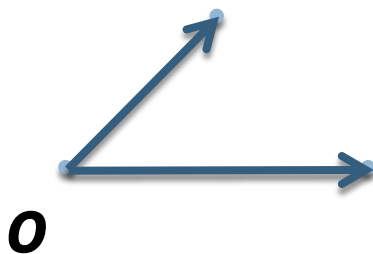


双对格子: 例

$$\mathbf{B} = \begin{pmatrix} 0 & 1 \\ 2 & 1 \end{pmatrix}$$

$$\mathbf{B}^{-T} = \begin{pmatrix} -1/2 & 1 \\ 1/2 & 0 \end{pmatrix}$$

$$\mathbf{B}^{-1} = \begin{pmatrix} -1/2 & 1/2 \\ 1 & 0 \end{pmatrix}$$

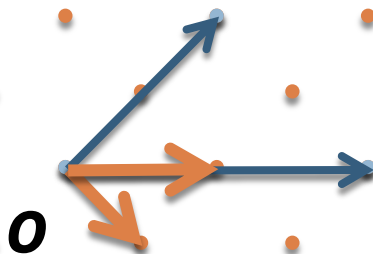


双对格子: 例

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$$\mathbf{B}^{-T} = \begin{pmatrix} -1/2 & 1 \\ 1/2 & 0 \end{pmatrix}$$

$$\mathbf{B}^{-1} = \begin{pmatrix} -1/2 & 1/2 \\ 1 & 0 \end{pmatrix}$$



特殊な格子

For $\mathbf{A} \in \mathbb{Z}_q^{n \times m}$

- $\Lambda_q^\perp(\mathbf{A}) = \{\mathbf{e} \in \mathbb{Z}^m \mid \mathbf{A}\mathbf{e} = \mathbf{0} \bmod q\}$
- $\Lambda_q(\mathbf{A}) = \{\mathbf{y} \in \mathbb{Z}^m \mid \mathbf{y} = \mathbf{A}^T \mathbf{s} \bmod q\}$

$\Lambda_q^\perp(\mathbf{A})$



$\Lambda_q(\mathbf{A})$



特殊な格子

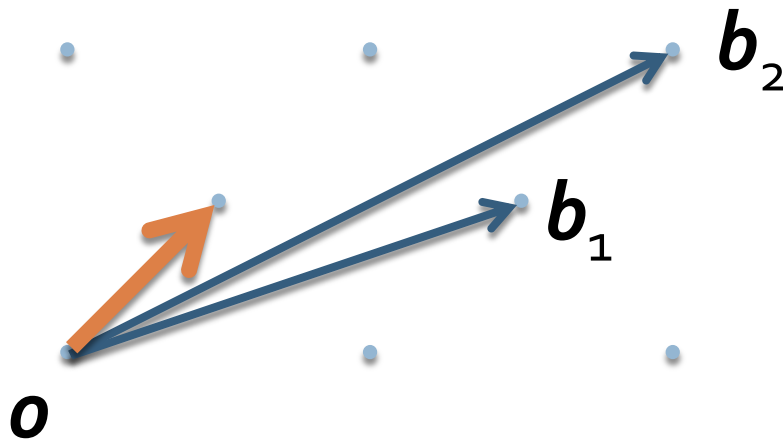
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$$(\Lambda_q(\mathbf{A}))^* = (1/q) \Lambda_q^\perp(\mathbf{A})$$

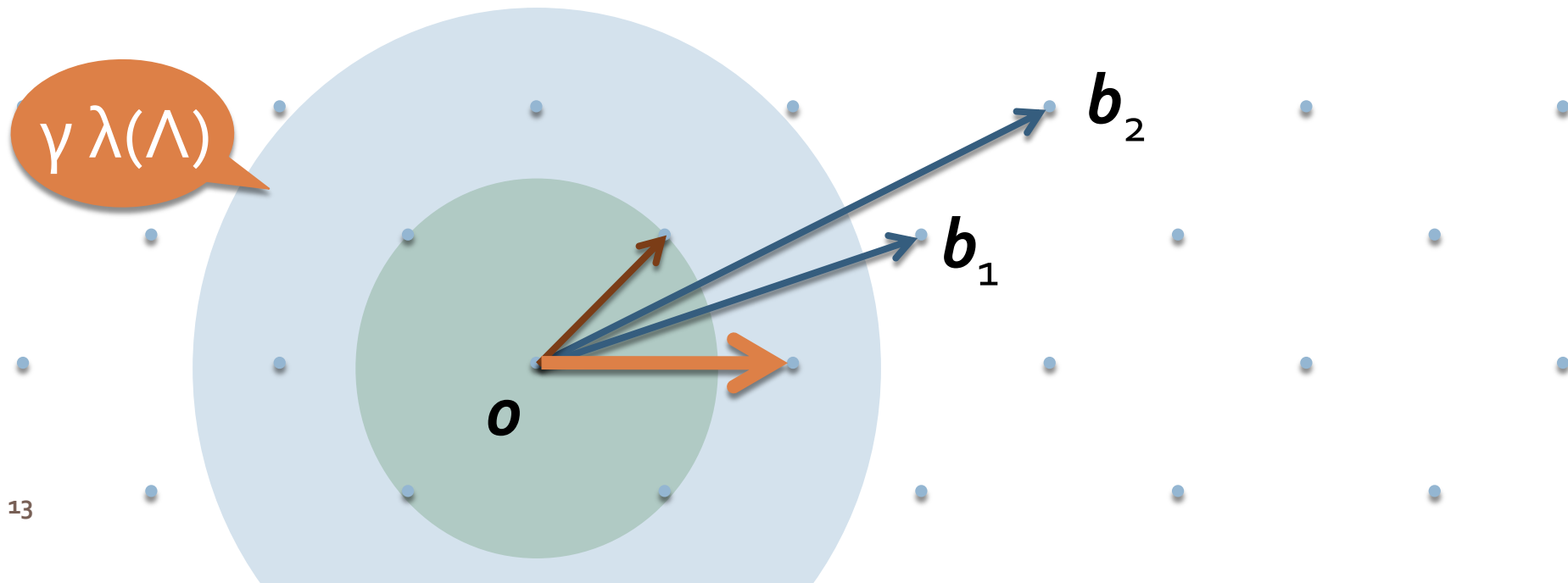
最短ベクトル問題 (SVP)

Given $\mathbf{B}$ of Λ
find $\mathbf{v} \in \Lambda - \{\mathbf{o}\}$ s.t. $\|\mathbf{v}\| = \lambda(\Lambda)$



近似版最短ベクトル問題 (SVP_γ)

Given $\mathbf{B}$ of Λ ,
find $\mathbf{v} \in \Lambda - \{\mathbf{o}\}$ s.t. $\|\mathbf{v}\| \leq \gamma \lambda(\Lambda)$



近似版最短線形独立ベクトル集合問題 (SIVP $_{\gamma}$)

Given $\mathbf{B}$ of Λ ,
find $\mathbf{S} \subseteq \Lambda$ s.t. $\|\mathbf{S}\| \leq \gamma \lambda_n(\Lambda)$

$\gamma \lambda_n(\Lambda)$

$\lambda_n(\Lambda)$

$\mathbf{o}$

$\mathbf{b}_2$

$\mathbf{b}_1$

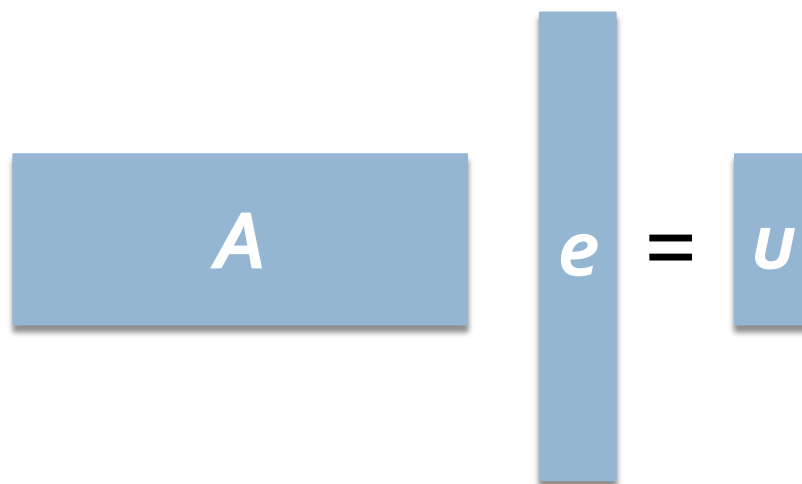
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イデ アル	✓	✓	✓	✓

格子ハッシュ [Ajt96,...,MR07,GPVo8]

$$H_n = \{h_A: D_n \rightarrow \mathbb{Z}_q^n \mid \mathbf{A} \in \mathbb{Z}_q^{n \times m}\}$$

$$h_A(\mathbf{e}) = \mathbf{A}\mathbf{e} \bmod q$$

$$D_n = \{\mathbf{e} \in \mathbb{Z}^m \mid \|\mathbf{e}\| \leq d\}$$



短い整数解問題 ($\text{SIS}_{q,m,\beta}$)

Given $\mathbf{A} \leftarrow \mathbb{Z}_q^{n \times m}$,
find $\mathbf{e} \in \mathbb{Z}^m$ s.t. $\|\mathbf{e}\| \leq \beta$, $\mathbf{A}\mathbf{e} = \mathbf{0} \bmod q$

格子 $\Lambda_q^\perp(\mathbf{A})$ の
短い非ゼロベクトルを
探索

$$\text{Col}(H) \geq \text{SIS}_{q,m,2d} \text{ [GGH96]}$$

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$$\text{Given } \mathbf{A} \leftarrow \mathbb{Z}_q^{n \times m},$$

$$\text{find } \mathbf{e} \in \mathbb{Z}^m \text{ s.t. } \|\mathbf{e}\| \leq \beta, \mathbf{A}\mathbf{e} = \mathbf{0} \bmod q$$

$$\text{SIS}_{q,m,\beta} \geq_{a/w} \text{SIVP}_{\tilde{O}(\beta\sqrt{n})} \quad [\text{Ajt96,GGH96,CNo7, Mico4,MRo7,GPVo8}]$$

Given $\mathbf{A} \leftarrow \mathbb{Z}_q^{n \times m}$,
 find $\mathbf{e} \in \mathbb{Z}^m$ s.t. $\|\mathbf{e}\| \leq \beta$, $\mathbf{A}\mathbf{e} = \mathbf{0} \bmod q$

Given $\mathbf{B}$ of Λ ,
 find $\mathbf{S} \subseteq \Lambda$ s.t. $\|\mathbf{S}\| \leq \gamma \lambda_n(\Lambda)$

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イデ アル	✓	✓	✓	✓

イデアル格子版の準備

$$R = \mathbb{Z}[x]/(1+x^4) \longleftrightarrow D \subseteq M_4(\mathbb{Z})$$

$$\mathbf{x}(x) = 1+x+x^3 \longleftrightarrow \text{Rot}_f(\mathbf{x}) = \begin{pmatrix} 1 & -1 & 0 & -1 \\ 1 & 1 & -1 & 0 \\ 0 & 1 & 1 & -1 \\ 1 & 0 & 1 & 1 \end{pmatrix}$$

$$\mathbf{y}(x) = 1+2x^2 \longleftrightarrow \text{Rot}_f(\mathbf{y}) = \begin{pmatrix} 0 & 0 & -2 & -1 \\ 1 & 0 & 0 & -2 \\ 2 & 1 & 0 & 0 \\ 0 & 2 & 1 & 0 \end{pmatrix}$$

$$\mathbf{x} \otimes \mathbf{y} = -1-3x+3x^2+2x^3 \longleftrightarrow \text{Rot}_f(\mathbf{x})\text{Rot}_f(\mathbf{y}) = \begin{pmatrix} -1 & -2 & -3 & 3 \\ -3 & -1 & -2 & -3 \\ 3 & -3 & -1 & -2 \\ 2 & 3 & -3 & -1 \end{pmatrix}$$

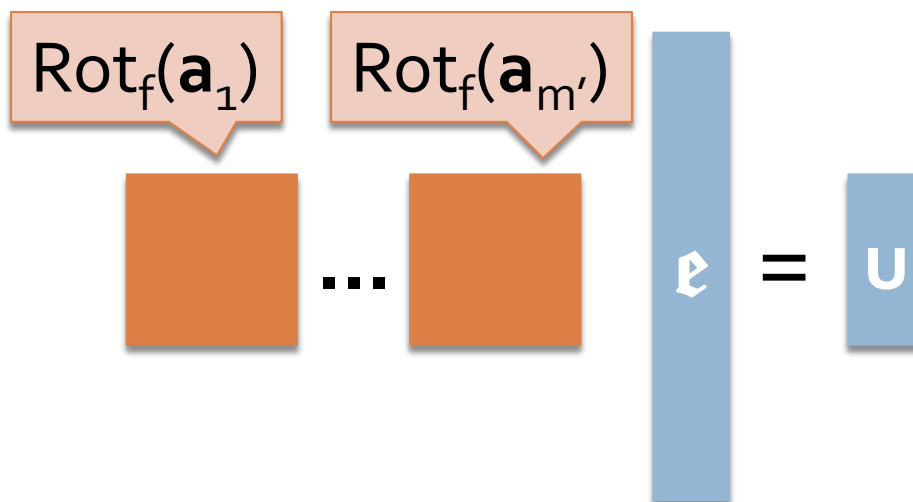
イデアル格子とイデアル格子問題

$$R = \mathbb{Z}[x]/\langle \mathbf{f} \rangle, \text{ e.g., } \mathbf{f} = x^n + 1$$
$$I \subseteq R \Leftrightarrow \Lambda_I \subseteq \mathbb{Z}^n$$

f-SVP_γ:
Given **B** of Λ_I ,
find $\mathbf{v} \in \Lambda_I - \{\mathbf{0}\}$ s.t. $\|\mathbf{v}\| \leq \gamma \lambda(\Lambda_I)$

イデアル格子ハッシュ [Mico7, LMo6, PRo6, PRo7]

$$\begin{aligned} H_n &= \{h_{\mathbf{a}}: D_n \rightarrow R_q \mid \mathbf{a} \in R_q^m\} \\ h_{\mathbf{a}}(\mathbf{e}) &= \mathbf{a} \mathbf{e} \bmod q = \sum_i \mathbf{a}_i \mathbf{e}_i \\ D_n &= \{\mathbf{e} \in R^m \mid \|\mathbf{e}\| \leq d\} \end{aligned}$$



$$\mathbf{f}\text{-Col}(H) \geq \mathbf{f}\text{-SIS}_{q,m,2d}$$

$$\begin{aligned} H_n &= \{h_{\mathbf{a}}: D_n \longrightarrow R_q \mid \mathbf{a} \in R_q^m\} \\ h_{\mathbf{a}}(\mathbf{e}) &= \mathbf{a} \mathbf{e} \bmod q = \sum_i \mathbf{a}_i \mathbf{e}_i \\ D_n &= \{\mathbf{e} \in R^m \mid \|\mathbf{e}\| \leq d\} \end{aligned}$$

Given $\mathbf{a} \leftarrow R_q^m$,
find $\mathbf{e} \in R^m$ s.t. $\|\mathbf{e}\| \leq \beta$, $\mathbf{a} \mathbf{e} = \mathbf{0} \bmod q$

$$\mathbf{f}\text{-SIS}_{q,m,\beta} \geq_{a/w} \mathbf{f}\text{-SVP}_{\tilde{O}(\beta\sqrt{n})} \quad [\text{Mico7, LMo6, PRo6, PRo7, SSTX09, LPR10}]$$

Given $\mathbf{a} \leftarrow R_q^m$,
 find $\mathbf{e} \in R^m$ s.t. $\|\mathbf{e}\| \leq \beta$, $\mathbf{a} \mathbf{e} = \mathbf{0} \bmod q$

$\mathbf{f}\text{-SVP}_\gamma$:
 Given $\mathbf{B}$ of Λ_I ,
 find $\mathbf{v} \in \Lambda_I - \{\mathbf{0}\}$ s.t. $\|\mathbf{v}\| \leq \gamma \lambda(\Lambda_I)$

$$\text{注: } \mathbf{f}\text{-SIS}_{q,m,\beta}^\infty \geq_{a/w} \mathbf{f}\text{-SVP}_{\tilde{O}(\beta)}^\infty$$

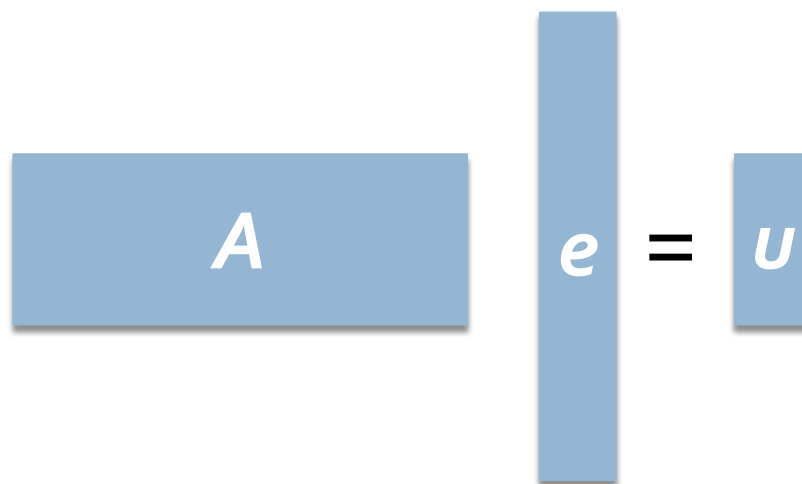
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	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

格子ハッシュ [Ajt96,...,MR07,GPVo8]

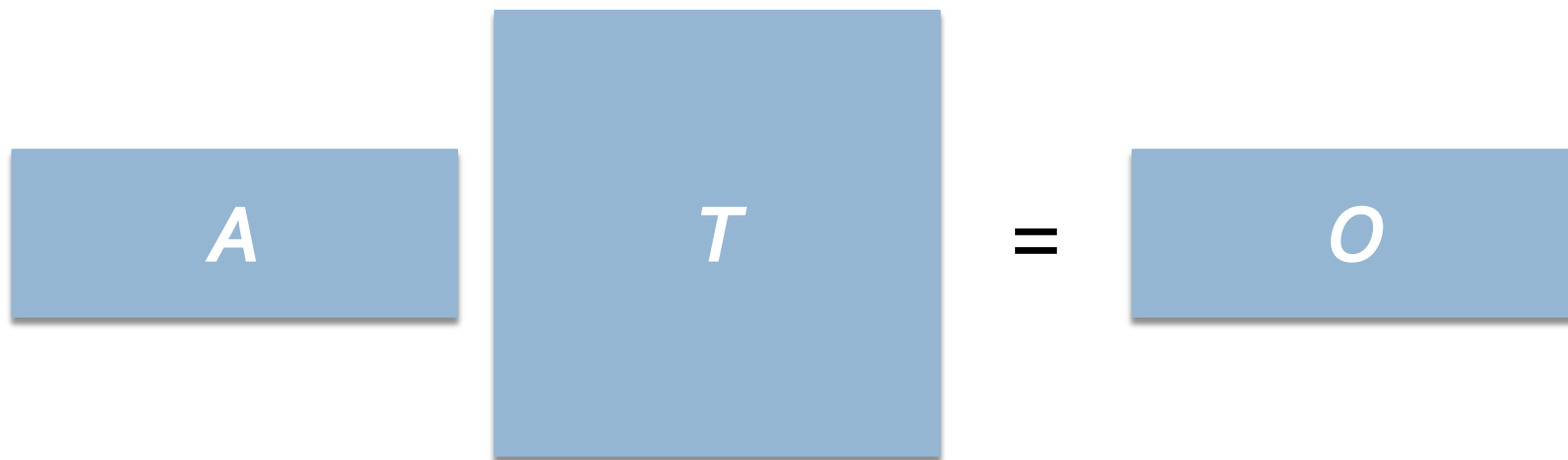
$$H_n = \{h_A: D_n \rightarrow \mathbb{Z}_q^n \mid \mathbf{A} \in \mathbb{Z}_q^{n \times m}\}$$

$$h_A(\mathbf{e}) = \mathbf{A}\mathbf{e} \bmod q$$

$$D_n = \{\mathbf{e} \in \mathbb{Z}^m \mid \|\mathbf{e}\| \leq d\}$$



落とし戸 [Ajt99, GPVo8, APo9]


$$A \cdot T = O$$

$$\begin{aligned} (\mathbf{A}, \mathbf{T}) &\leftarrow \text{Gen}(1^n) \\ \mathbf{T} &\in \mathbb{Z}^{m \times m} \text{ s.t.} \\ \mathbf{T} &\text{: a basis of } \Lambda_q^\perp(\mathbf{A}), \|\text{GS}(\mathbf{T})\| \leq L \end{aligned}$$

注: $L = \tilde{O}(\sqrt{n \log q})$

落とし戸 [Ajt99, GPVo8, AP09]

$$H_n = \{h_A: D_n \rightarrow \mathbb{Z}_q^n \mid \mathbf{A} \in \mathbb{Z}_q^{n \times m}\}$$

$$h_A(\mathbf{e}) = \mathbf{A}\mathbf{e} \bmod q$$

$$D_n = \{\mathbf{e} \in \mathbb{Z}^m \mid \|\mathbf{e}\| \leq d\}$$

τ があれば,
 $h_A(\mathbf{e}) = \mathbf{u}$ になる $\mathbf{e}$ を取れる
(ただし $\|\mathbf{e}\| \leq L\sqrt{m} \cdot \omega(\sqrt{\log m})$)

$$L = O(\sqrt{m \log q})$$

ガウス分布

- ガウス分布関数:

$$v_{s,c}(\mathbf{x}) = \exp(-\pi \|(\mathbf{x}-\mathbf{c})/s\|^2) / s^n$$

- 離散ガウス分布関数:

$$D_{\Lambda,s,c}(\mathbf{x}) = v_{s,c}(\mathbf{x}) / v_{s,c}(\Lambda)$$

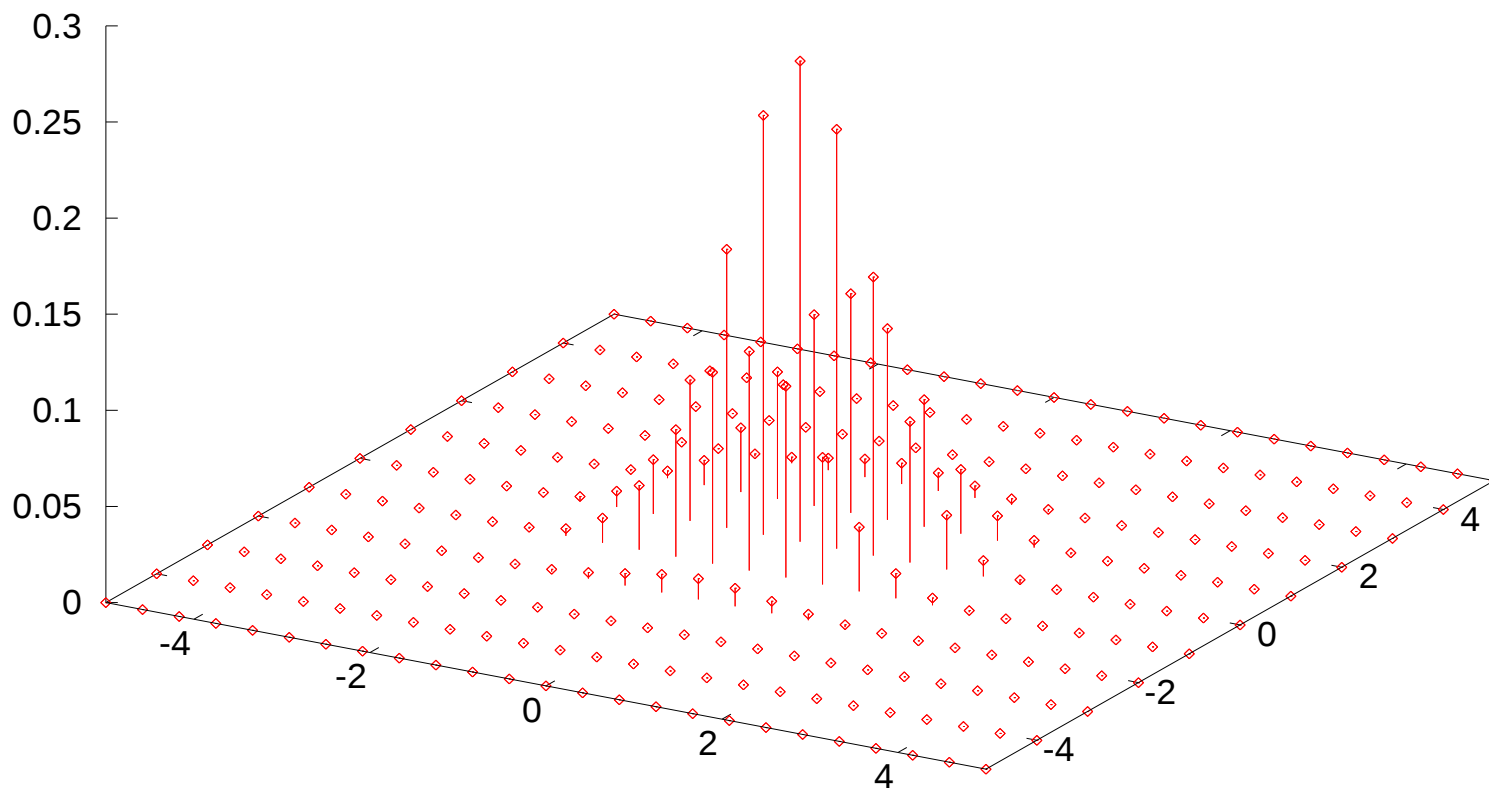
- $s \geq \lambda_n(\Lambda) \omega(\sqrt{\log n})$ なら,

$$\Pr_{\mathbf{x} \leftarrow D}[\|\mathbf{x}\| > s\sqrt{n}] < \text{negl}(n)$$

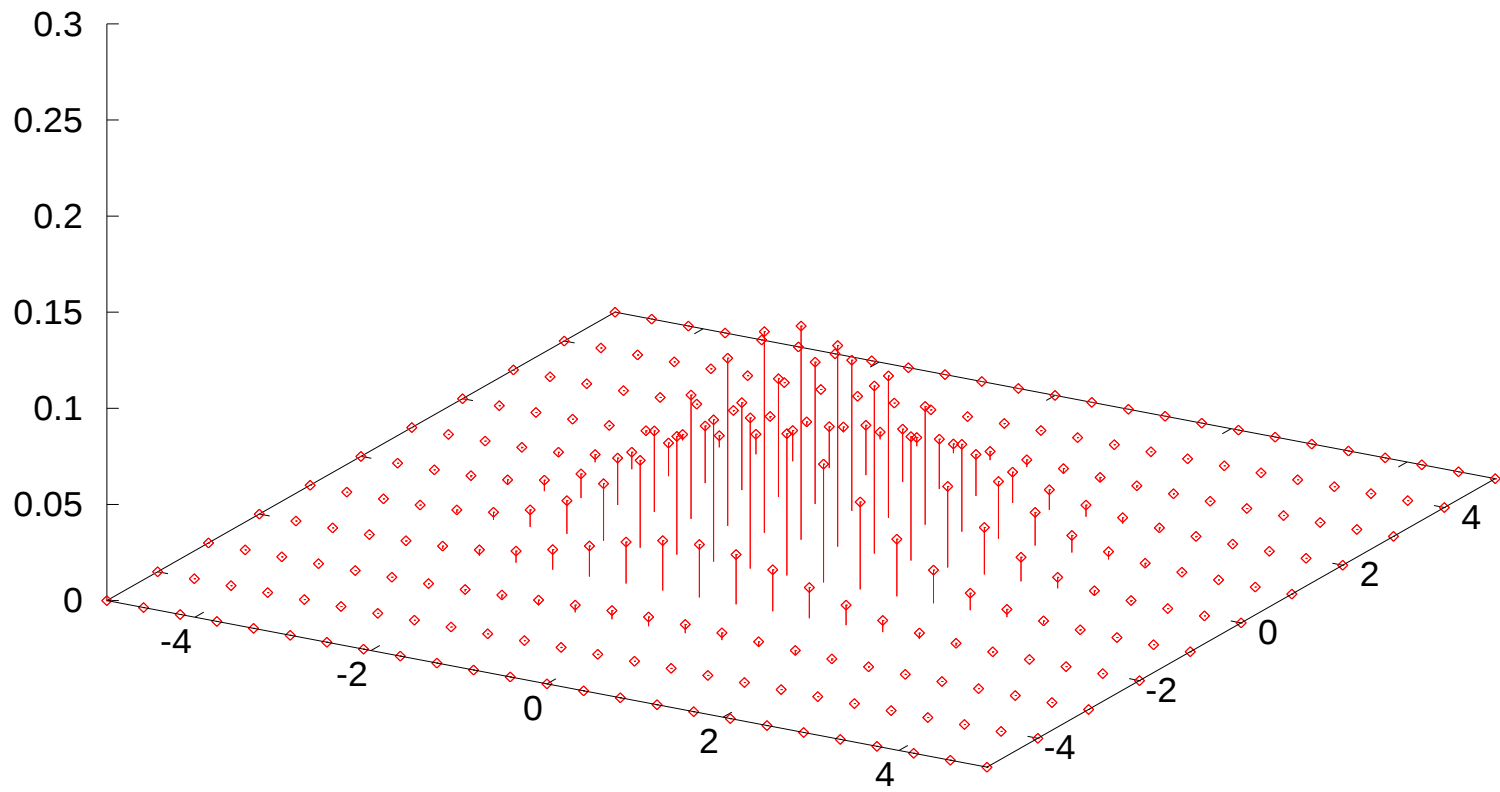
図の気持ち

- ガウス分布の性質を保っていそう
- s のサイズが小さいとガウス分布の性質を保たなさそう

$D_{L,2}$



$$D_{L,3}$$



Klein-GPV Sampling [Kle01?, GPVo8, Pei10]

$\mathbf{T}$: a basis of Λ , $\|\text{GS}(\mathbf{T})\| \leq L$,
 $s \geq L \omega(\sqrt{\log n})$,
 $\Rightarrow \text{SampleD}(\mathbf{T}, s, \mathbf{c}) \sim_S D_{\Lambda, s, \mathbf{c}}$

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$(\mathbf{A}, \mathbf{T}) \leftarrow \text{Gen}(1^n)$
 $\mathbf{T} \in \mathbb{Z}^{m \times m}$ s.t.
 $\mathbf{T}$: a basis of $\Lambda_q^\perp(\mathbf{A})$, $\|\text{GS}(\mathbf{T})\| \leq L$

注: $L = O(\sqrt{n \log q})$

Klein-GPV Sampling [Kle01?, GPVo8, Pei10]

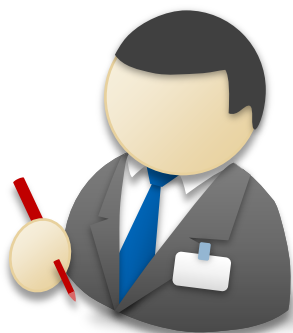
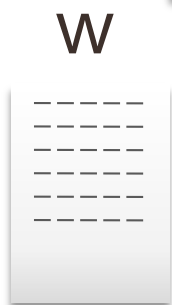
$\mathbf{T}$: a basis of Λ , $\|\text{GS}(\mathbf{T})\| \leq L$,
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 $\Rightarrow \text{SampleD}(\mathbf{T}, s, \mathbf{c}) \sim_S D_{\Lambda, s, \mathbf{c}}$

$\mathbf{T}$ があれば,
 $h_A(\mathbf{e}) = \mathbf{u}$ になる $\mathbf{e}$ を取れる
 (ただし $\|\mathbf{e}\| \leq s\sqrt{m}$, $s = L \omega(\sqrt{\log m})$)

GPV署名 [GPVo8]

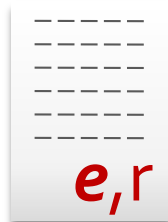
1. $u = H(w || r)$
2. $e \leftarrow h_A^{-1}(u)$

1. $h_A(e) = H(w || r)$
2. $\|e\| \leq s\sqrt{m}$?



$vk = A$
 $sk = T$

$w, (e, r)$



$vk = A$

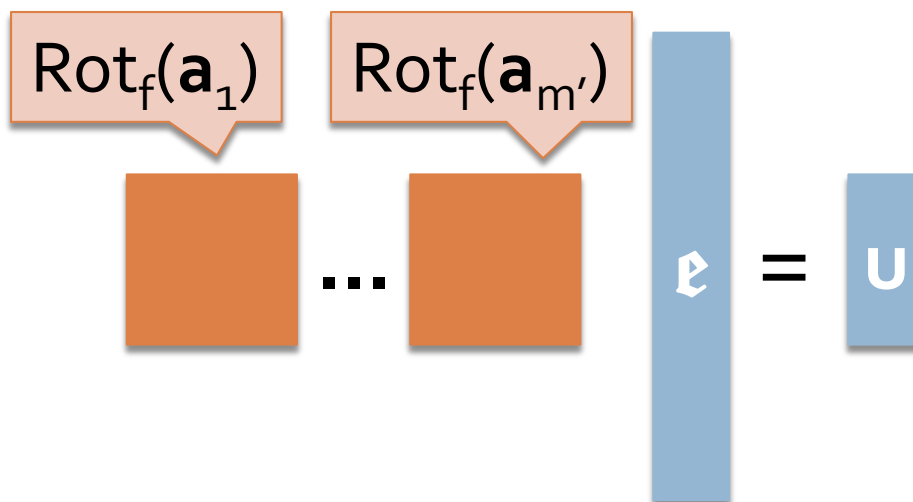
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イデ アル	✓	✓	✓	✓

イデアル格子ハッシュ [Mico7, LMo6, PRo6, PRo7]

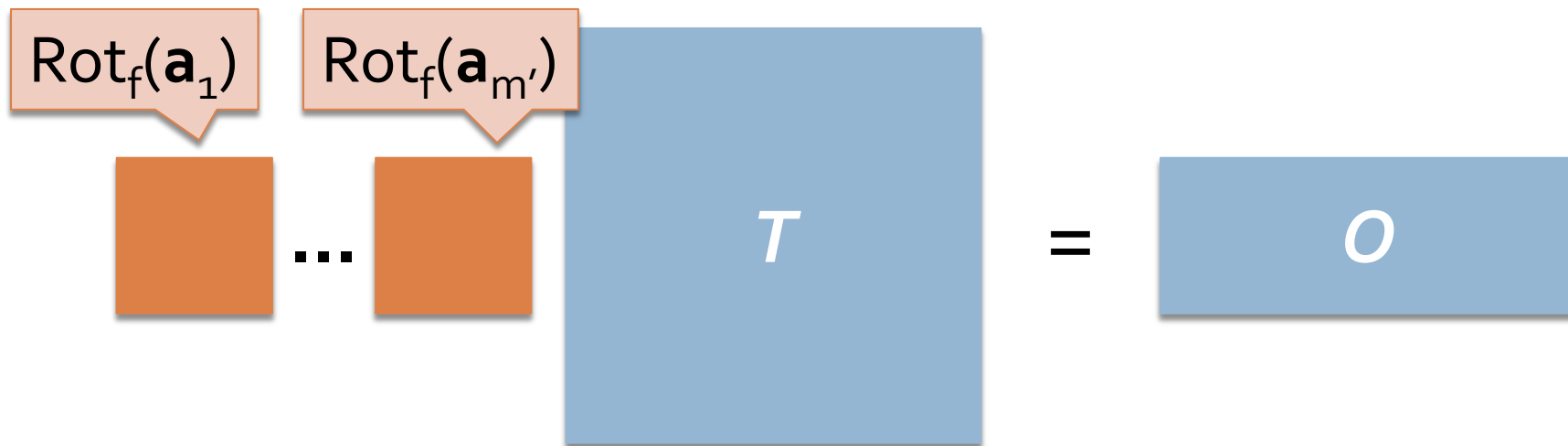
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$$h_{\mathbf{a}}(\mathbf{e}) = \mathbf{a} \mathbf{e} \bmod q = \sum_i \mathbf{a}_i \mathbf{e}_i$$

$$D_n = \{\mathbf{e} \in R^m \mid \|\mathbf{e}\| \leq d\}$$



落とし戸 [SSTX09]



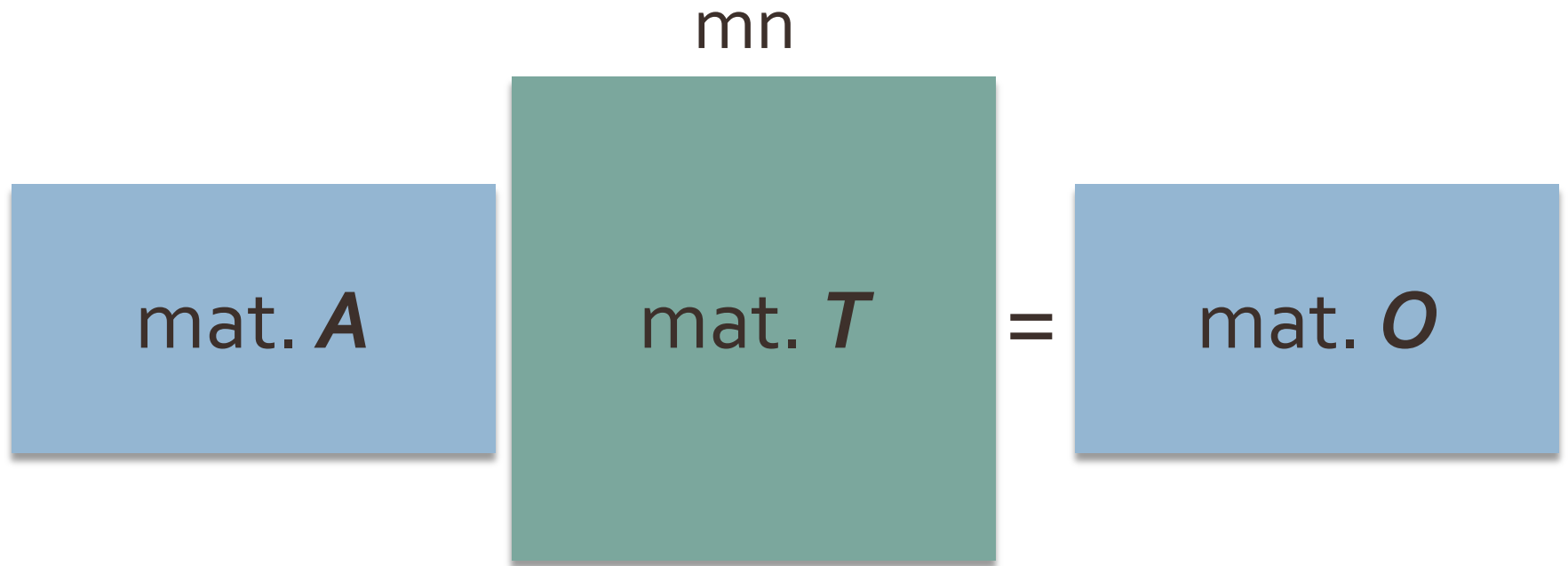
$$(\mathbf{a} \in \mathbb{R}_q^m, T) \leftarrow \text{Gen}(1^n)$$

$$T \in \mathbb{Z}^{mn \times mn} \text{ s.t.}$$

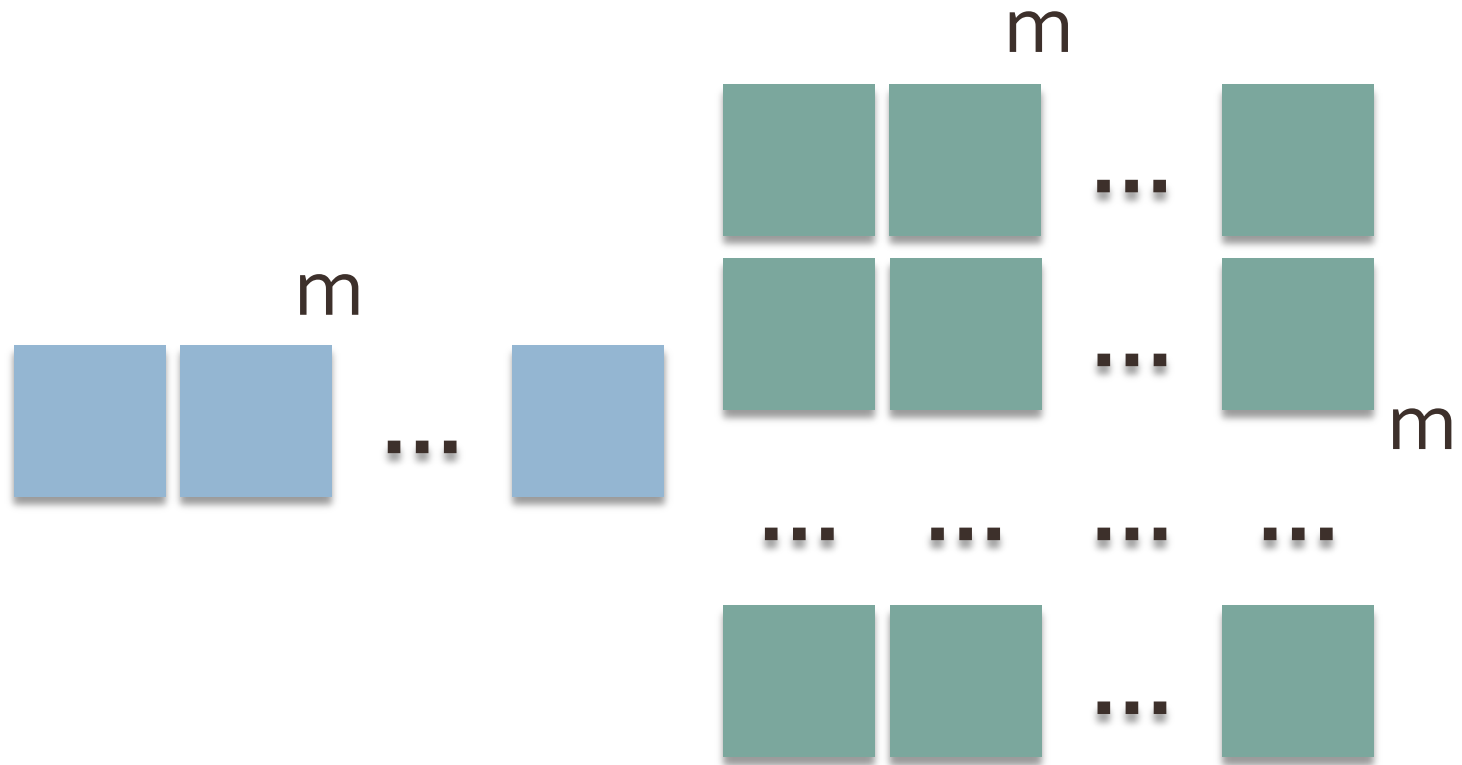
$$T: \text{a basis of } \Lambda_q^\perp(\text{Rot}(\mathbf{a})), \|\text{GS}(T)\| \leq L$$

注: $m = O(\log^2 q)$ で $L = O(\sqrt{n \log q})$
 $m = O(\log q)$ で $L = \tilde{O}(n \log q)$

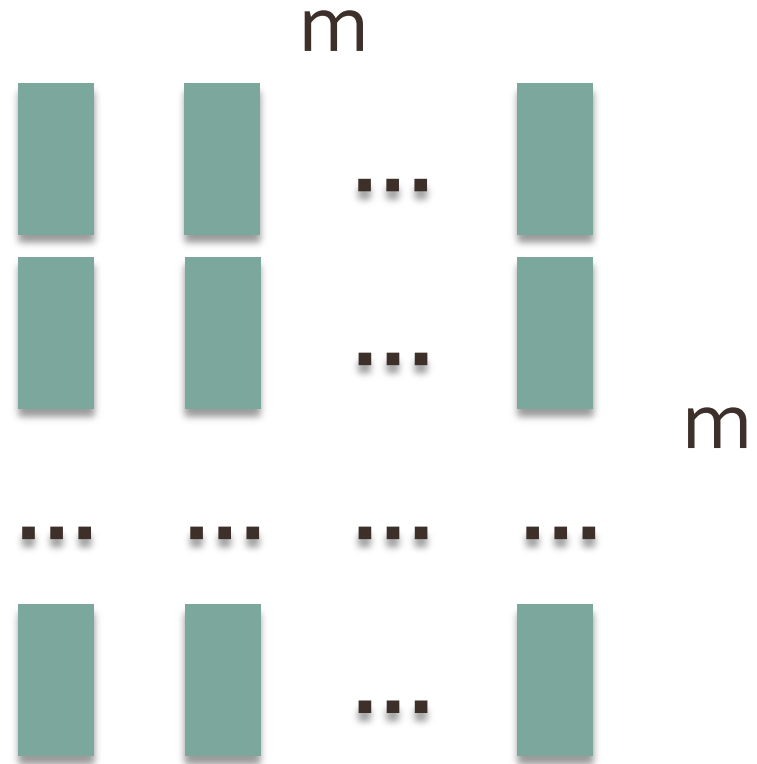
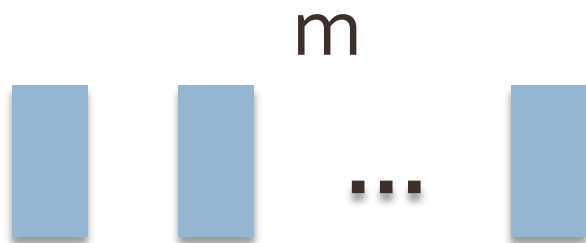
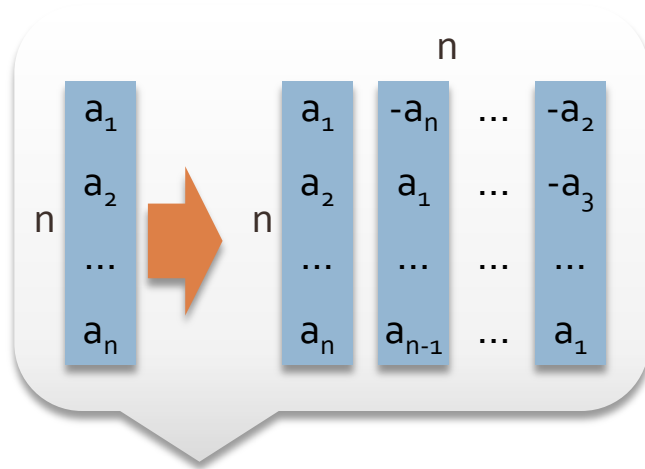
イデアル版鍵生成アルゴリズム [SSTX09]



イデアル版鍵生成アルゴリズム [SSTX09]



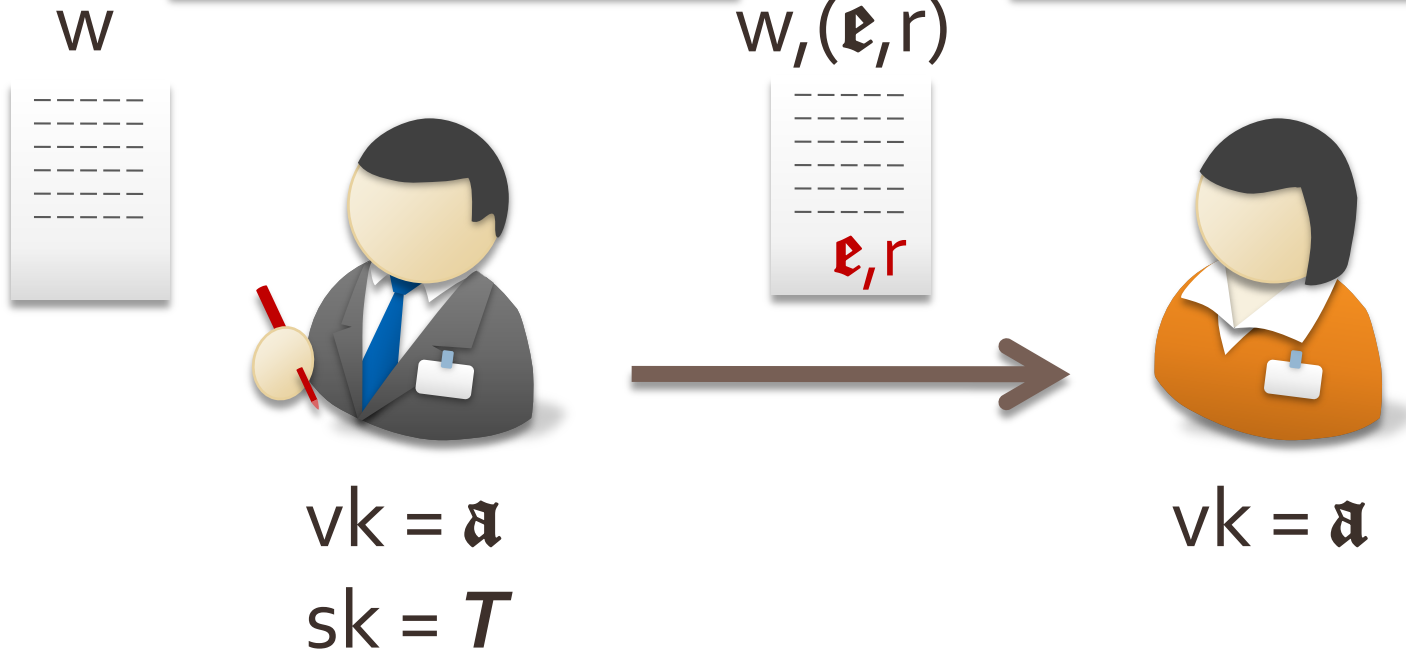
イデアル版鍵生成アルゴリズム [SSTX09]



SSTX署名 [SSTX09]

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1. $h_a(e) = H(w || r)$
2. $\|e\| \leq s\sqrt{mn}$?



	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MRo7, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

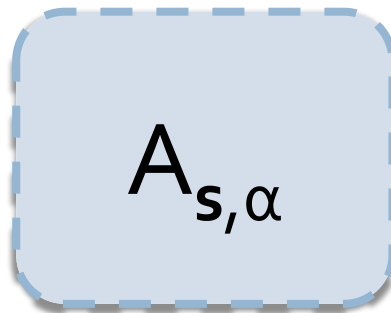
LWE関数 [Regog,GPVo8,Peiog]

$$\mathbf{A} \in \mathbb{Z}_q^{n \times m}, \mathbf{s} \leftarrow \mathbb{Z}_q^n, \mathbf{x} \leftarrow \chi_\alpha^m$$
$$g_A(\mathbf{s}, \mathbf{x}) = \mathbf{A}^T \mathbf{s} + \mathbf{x} \bmod q$$

$$\mathbf{A}^T \mathbf{s} + \mathbf{x} = \mathbf{p}$$

sLWE_{q,α}

Find $\mathbf{s} \leftarrow \mathbb{Z}_q^n$
given $A_{\mathbf{s},\alpha} \rightarrow (\mathbf{a}, \langle \mathbf{a}, \mathbf{s} \rangle + x)$



Find $\mathbf{s}$

$A_{\mathbf{s},\alpha}$
1. $\mathbf{a} \leftarrow \mathbb{Z}_q^n, x \leftarrow \chi_\alpha$
2. $b = \langle \mathbf{a}, \mathbf{s} \rangle + x$
3. Output $(\mathbf{a}, b)$

落とし戸 [Peio9?]

$$\mathbf{A} \in \mathbb{Z}_q^{n \times m}, \mathbf{s} \leftarrow \mathbb{Z}_q^n, \mathbf{x} \leftarrow \chi_\alpha^m$$
$$g_A(\mathbf{s}, \mathbf{x}) = \mathbf{A}^T \mathbf{s} + \mathbf{x} \bmod q$$
$$\mathbf{T}: \text{a basis of } \Lambda_q^\perp(\mathbf{A})$$

$$\mathbf{d} = \mathbf{T}^T \mathbf{p} \bmod q (= \mathbf{T}^T \mathbf{x} \bmod q)$$
$$\mathbf{c} = \mathbf{T}^{-T} \mathbf{d} \text{ (in } \mathbb{Q}) (= \mathbf{x} \text{ with ow. prob)}$$

Extract $\mathbf{s}$ from $\mathbf{v} = \mathbf{p} - \mathbf{c} = (\mathbf{A}^T \mathbf{s} \bmod q)$

注: $|\langle \mathbf{t}_i, \mathbf{x} \rangle| \leq L a \leq q/2 \Rightarrow a \leq q/2L$ なら $\mathbf{c}$ の行が成立
 $q \geq 5L\sqrt{m}$, $1/\alpha \leq L \omega(\sqrt{\log n})$ とする

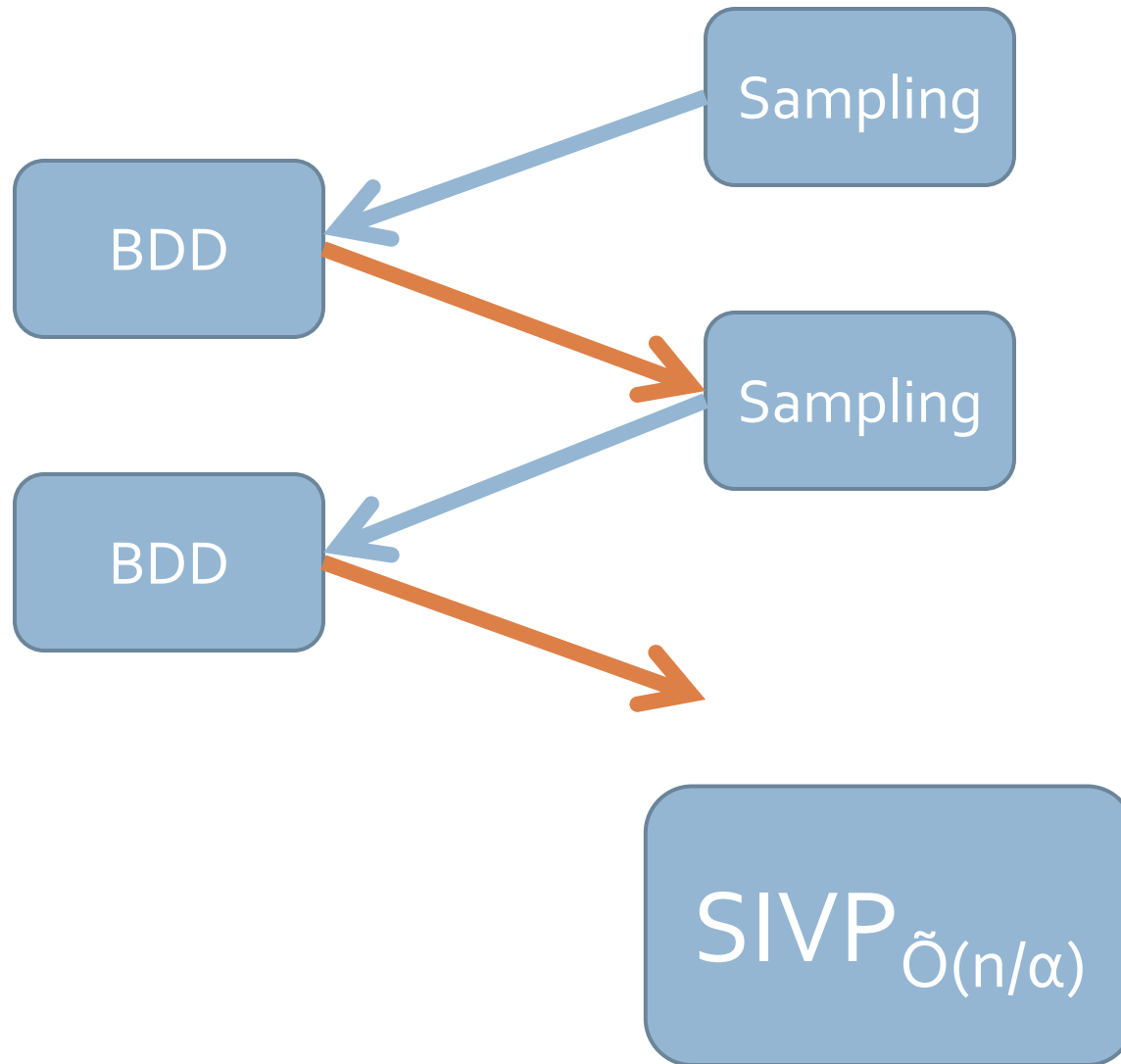
追記: LWE関数 [Regog, GPVo8, Peiog]

$$\mathbf{A} \in \mathbb{Z}_q^{n \times m}, \mathbf{s} \leftarrow \mathbb{Z}_q^n, \mathbf{x} \leftarrow \chi_\alpha^m$$
$$g_A(\mathbf{s}, \mathbf{x}) = \mathbf{A}^\top \mathbf{s} + \mathbf{x} \bmod q$$

$$(\mathbf{A}, p) \sim_c (\mathbf{A}, U(\mathbb{Z}_q^m))$$

も言える

Regevの帰着 [Reg09]



	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MRo7, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

イデアル格子版 $g_{\mathfrak{a}}(s, x)$

$$\mathfrak{a} \in R_q^m, s \leftarrow \mathbb{Z}_q^n, x \leftarrow \chi_\alpha^{mn}$$

$$g_{\mathfrak{a}}(s, x) = A^T s + x \bmod q$$

$$A = \text{Rot}_f(\mathfrak{a}) = [\text{Rot}_f(\mathbf{a}_1), \dots, \text{Rot}_f(\mathbf{a}_m)]$$

The diagram illustrates the decomposition of the matrix A^T into its constituent vectors. It shows two equivalent equations for the ideal lattice function $g_{\mathfrak{a}}(s, x)$.

On the left, the full equation is shown: a large blue rectangle labeled A^T is multiplied by a small blue rectangle labeled s , followed by a plus sign and a large blue rectangle labeled x , an equals sign, and a large blue rectangle labeled p .

On the right, the same equation is shown, but the matrix A^T is decomposed into m rows. Each row is represented by a small blue rectangle labeled s multiplied by a large orange rectangle. The top orange rectangle is labeled $\text{Rot}_f(\mathbf{a}_1)^T$ and the bottom one is labeled $\text{Rot}_f(\mathbf{a}_m)^T$. This is followed by a plus sign and a large blue rectangle labeled x , an equals sign, and a large blue rectangle labeled p .

f-sLWE_{q,α}

$\text{Rot}_f(\mathbf{a}_1)^\top$

$\mathbf{s}$

+

$\mathbf{x}$

Find $\mathbf{s} \leftarrow \mathbb{Z}_q^n$,
given $A_{\mathbf{s},\alpha} \rightarrow (\mathbf{a}, \text{Rot}_f(\mathbf{a})^\top \mathbf{s} + \mathbf{x})$



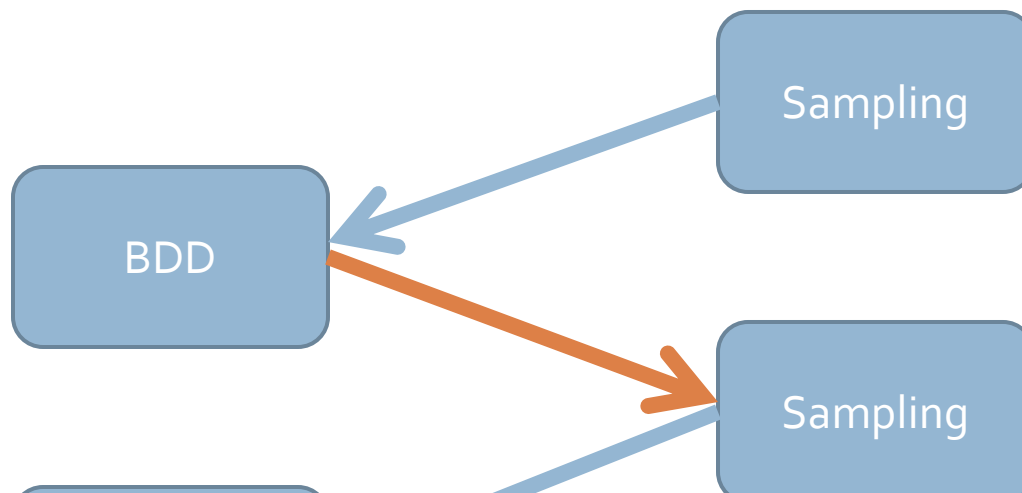
$A_{\mathbf{s},\alpha}$

Find $\mathbf{s}$

$A_{\mathbf{s},\alpha}$

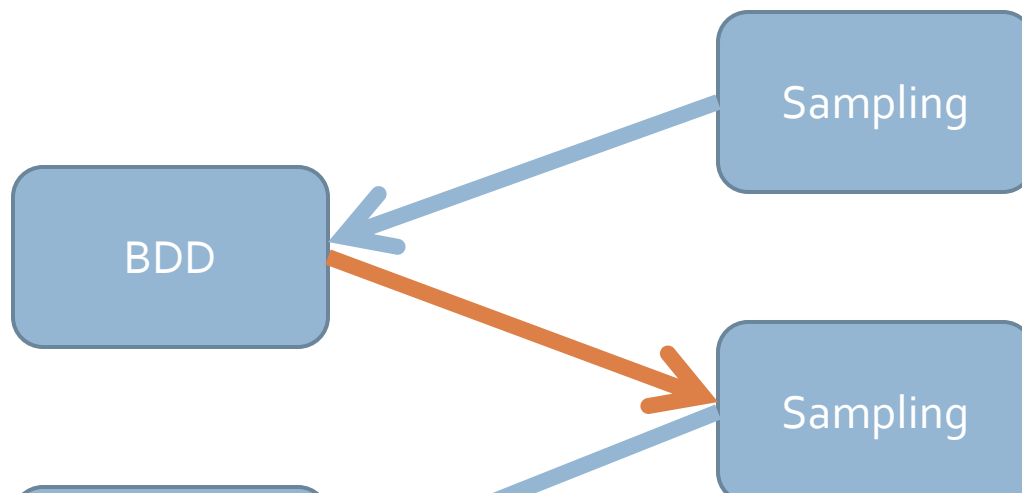
1. $\mathbf{a} \leftarrow R_q, \mathbf{x} \leftarrow \chi_\alpha^n$
2. $\mathbf{b} = \text{Rot}_f(\mathbf{a})^\top \mathbf{s} + \mathbf{x}$
3. Output $(\mathbf{a}, \mathbf{b})$

Regevの帰着



sLWE+古典帰着が
イデアル格子だと通らない

Regevの帰着



sLWE+古典帰着を諦めて
量子帰着だけ使えばいいよ

[Regog] の再解釈

量子帰着 w/w

sLWE+古典帰着 w/w



Given $\mathbf{B}$ of Λ^* ,
If $R(\Lambda, \mathbf{p})$ finds $\mathbf{v}$ in Λ s.t. $\|\mathbf{v} - \mathbf{p}\| \leq d$,
 $S^R(\mathbf{B})$ samples from $D_{\Lambda^*, s}$



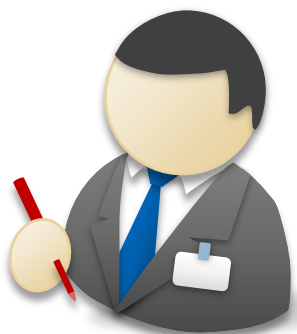
$\Lambda = \Lambda_q(\mathbf{A})$ とすると sLWE で $\mathbf{A}^T \mathbf{s}$ が分かる
 $(\Lambda_q(\mathbf{A}))^* = (1/q) \Lambda_q^\perp(\mathbf{A})$ の短いベクトルが分かる
 $\Rightarrow \text{SIS}_{q, m, \beta}$ が解ける

	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MRo7, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

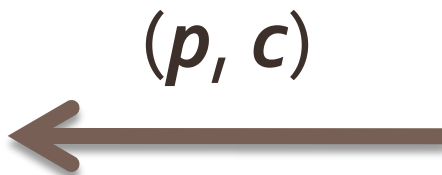
LWE-TDFs based KEM [Peio9]

1. $s \leftarrow p$
2. $k \leftarrow c - v$

1. $p = g_A(s, x)$
2. $v = g_U(s, x')$
3. $c = v + kq/2$



$ek = (A, U)$
 $dk = T$



$ek = (A, U)$

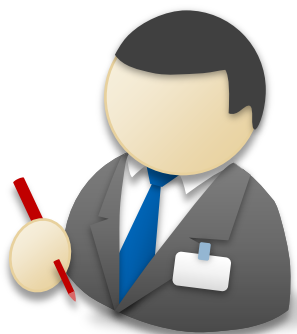
$(A, p) \sim_c (A, U(\mathbb{Z}_q^m))$ を利用

	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MRo7, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓

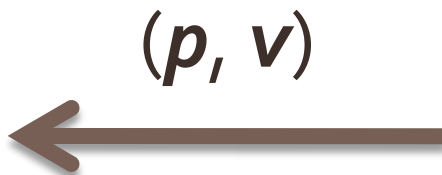
f-sLWE-TDFs based PKE [SSTX09]

1. $s \leftarrow p$
2. $w \leftarrow v \oplus H(s)$

1. $p = g_a(s, x)$
2. $v = w \oplus H(s)$



$ek = a$
 $dk = T$



$ek = a$

注: $(a, p) \sim_c (a, U(\mathbb{Z}_q^m))$ は不明

まとめ－色々やった

	Hash	ID	CR-PSFs	Sig (w/ ROM)
一般	[Ajt97,...,MR07, GPVo8]	[MVo3,Lyu08, KTXo8]	[GPVo8]	[GPVo8]
イデ アル	[Mico4,LMo6, PRo6]	[Lyu08,KTXo8, Lyu09]	✓	✓
	OW-TDFs	PKE, IBE HIBE	IBI	Sig (w/o ROM)
一般	[GPVo8,Peiog]	[Rego5,GPVo8, Peiog,CHKP10]	✓	[CHKP10]
イデ アル	✓	✓	✓	✓